

ECE 510 OCE

BDDs and Their Applications

Lecture 18.

Verification Using Specialized DDs

May 30, 2000
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Overview

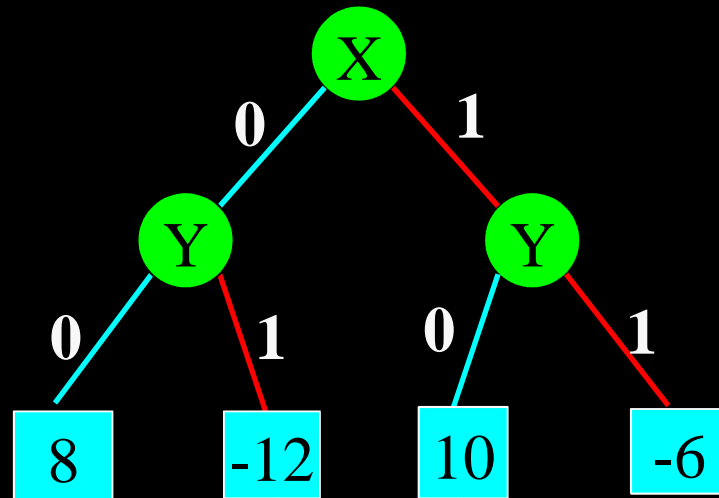
- Extensions of Binary Decision Diagrams
 - Multi-terminal BDDs (MTBDDs) (Clarke, DAC'93)
 - (Multiplicative) binary moment diagrams (BMDs/*BMDs) (Bryant, DAC'95)
 - Edge-valued BDDs (EVBDDs) (Lai, DAC'92)
 - Kronecker functional decision diagrams (OKFDDs) (Drechsler/Sarabi/Perkowski, DAC'94)
 - Kronecker multiplicative moment diagrams (K*BMDs) (Drechsler, EDTC'96)
- Word-level DDs vs. bit level DDs for verification
- Binary expression diagrams (Andersen/Hulgaard, 1997)
- Advantages and limitations of the specialized DDs

Extensions of BDDs

- Representation of “pseudo-Boolean” functions: integer (real) number functions over Boolean functions
- **Applications:** integer linear programming, matrix multiplication, spectral transforms, word-level analysis of digital systems

Multi Terminal BDDs

X	Y	F
0	0	8
0	1	-12
1	0	10
1	1	-6



Generalization of Expansion

- Boole-Shannon expansion:

$$F = x' \& F_0 + x \& F_1$$

- For numeric value functions, generalized to

$$F = (1-x) \& F_0 + x \& F_1$$

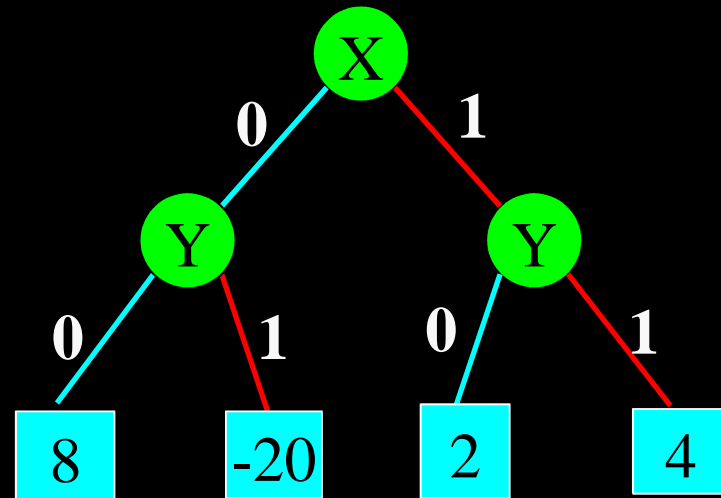
- Boole-Shannon expansion can be rewritten as

$$F = F_0 + x \& (F_1 - F_0)$$

where $F_x = F_1 - F_0$ is the **binary moment**
(derivative of F w.r.t. variable x)

Binary Moment Diagrams (BMDs)

X	Y	F
0	0	8
0	1	-12
1	0	10
1	1	-6



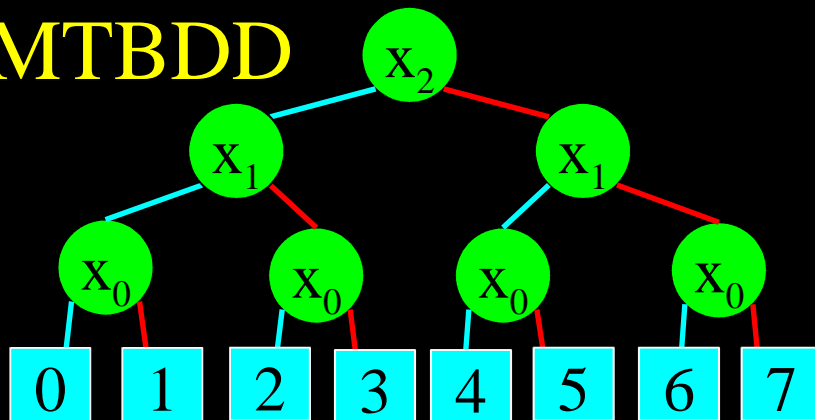
$$\begin{aligned} F &= 8(1-x)(1-y) - 12(1-x)y + 10x(1-y) - 6xy = \\ &= 8 - 20y + 2x + 4xy \end{aligned}$$

Edge-Valued BDDs

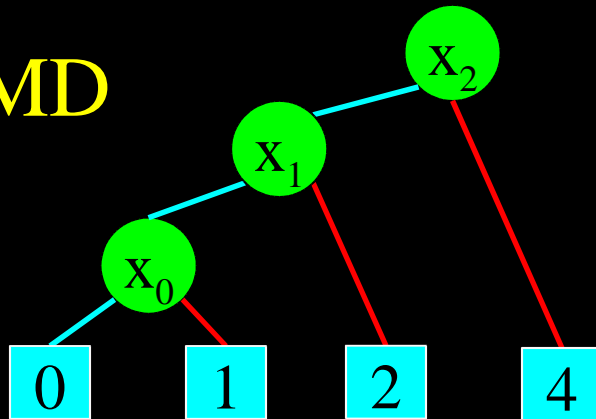
- An edge has an integer weight
- The weights are combined **additively**
- The value of the function is derived by following a path from the root to a leaf and **summing** the edge weights encountered

Example: $F = X = 4x_2 + 2x_1 + x_0$

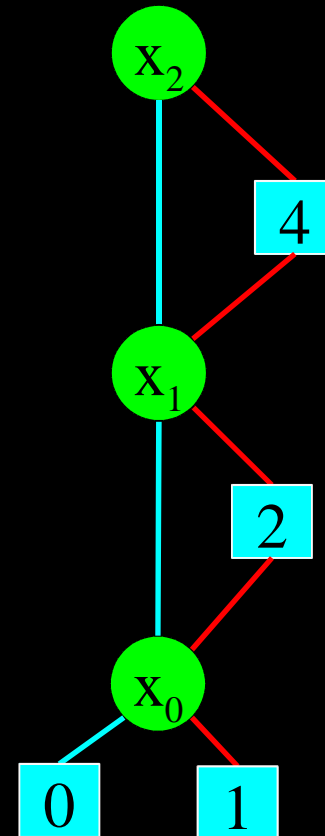
MTBDD



BMD



EVBD

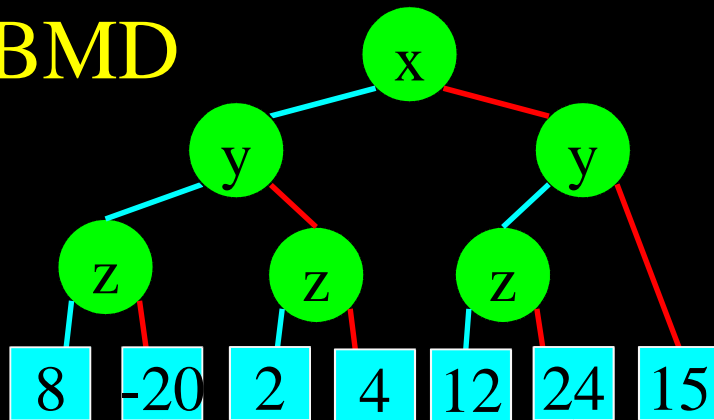


Multiplicative BMDs

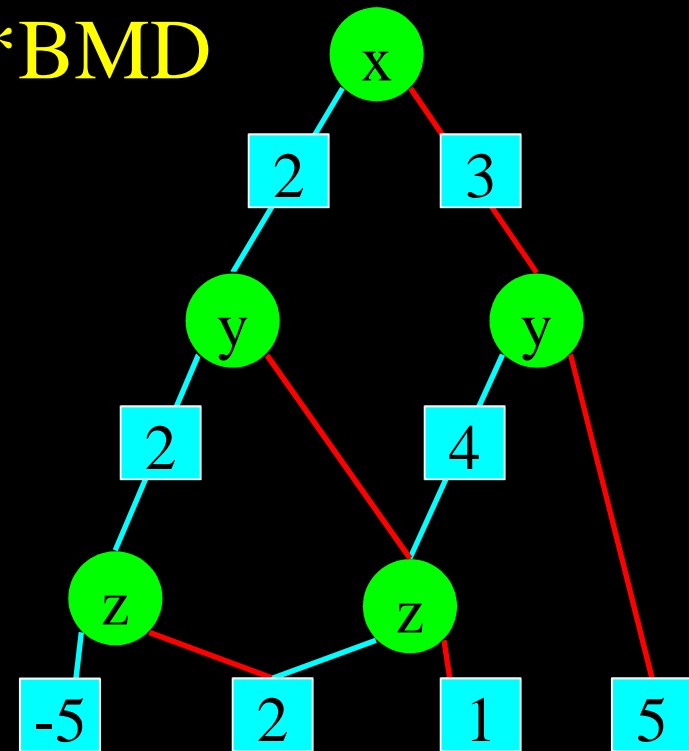
- An edge has an integer (rational) weight
- The weights are combined **multiplicatively**
- The value of the function is derived by following a path from the root to a leaf and **multiplying** the edge weights encountered

Example: $F=8-20z+2y+4yz+12x+24xz+15xy$

BMD



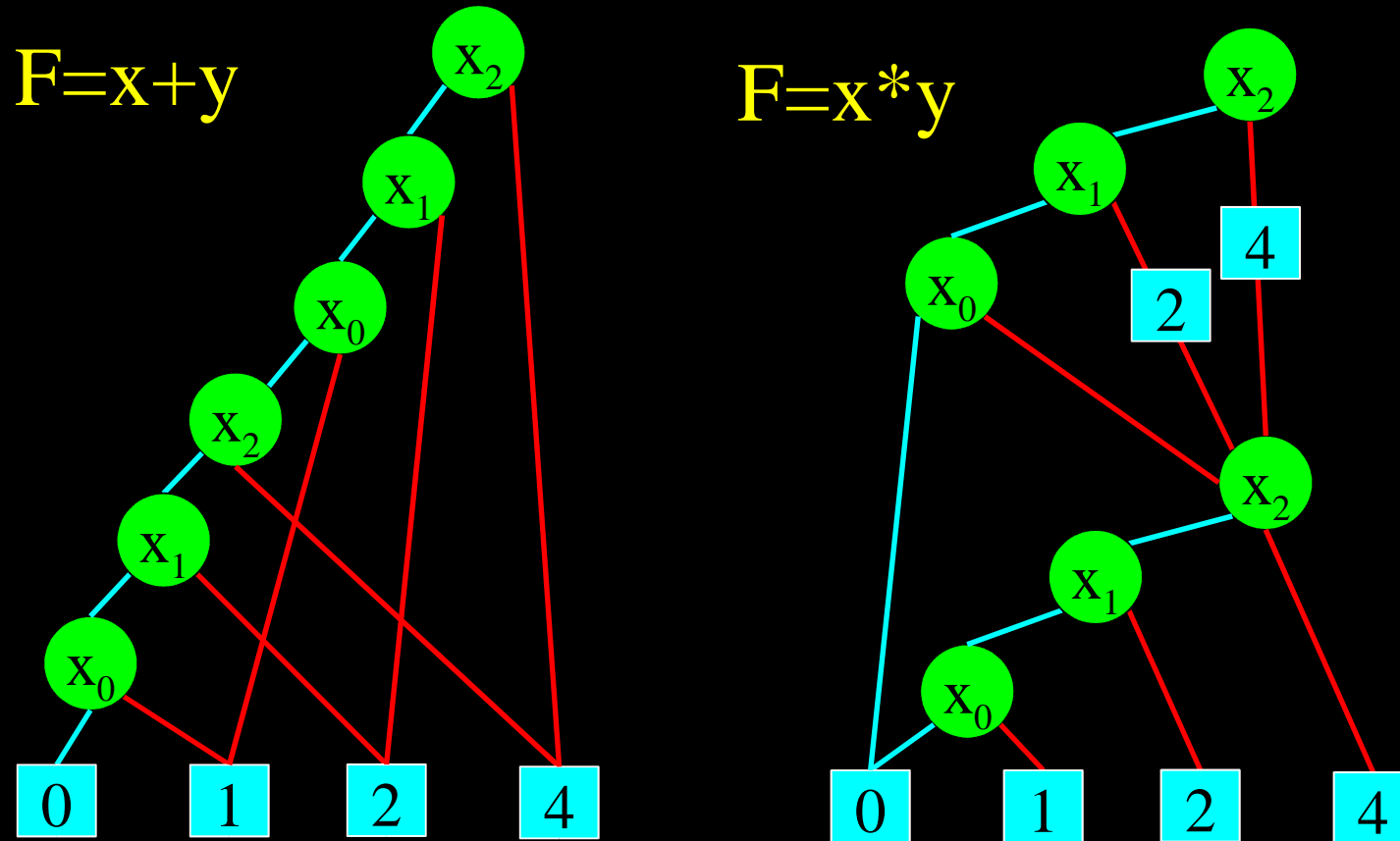
*BMD



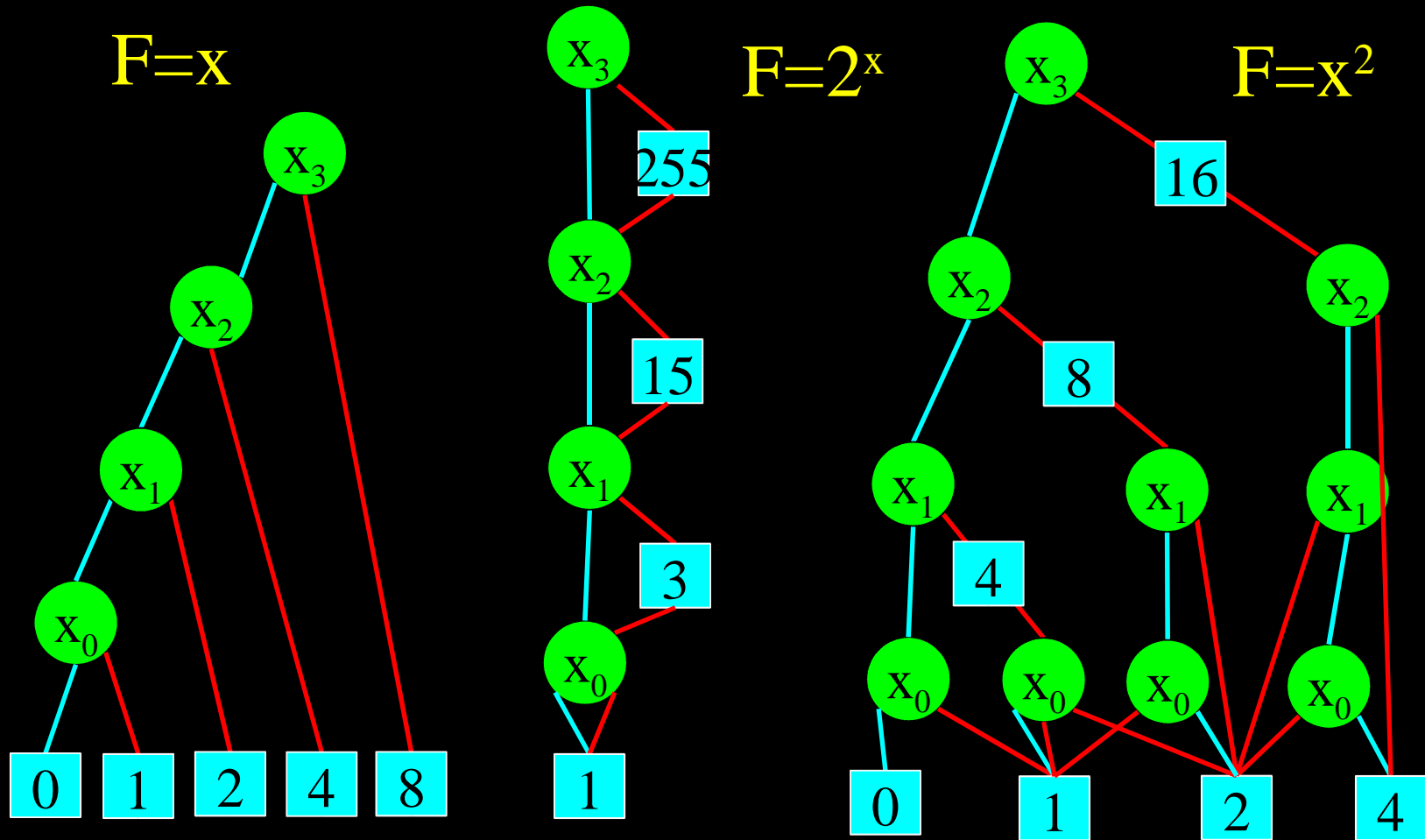
Word-Level Operation Complexity

Form	x	$x+y$	$x*y$	x^2	c^x
MTBDD	exp	exp	exp	exp	exp
EVBDD	lin	lin	exp	exp	exp
BMD	lin	lin	quadr	quadr	exp
*BMD	lin	lin	lin	quadr	lin
K*BMD	lin	lin	lin	quadr	lin

Sum/Product *BMD Representation

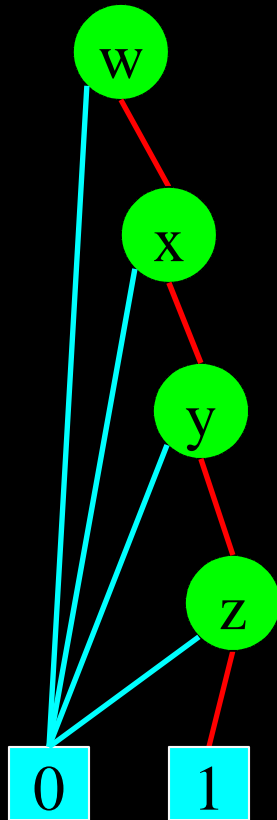


Unary Word-Level *BMD Operations

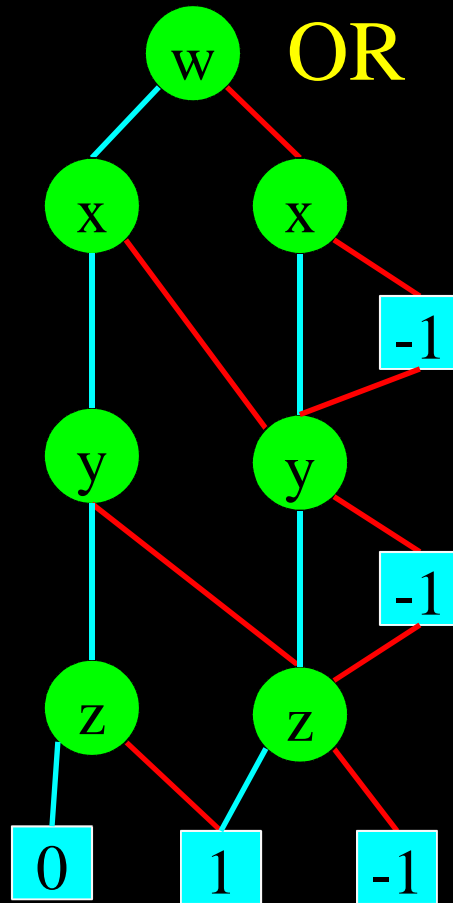


*BMDs for Common Boolean Functions

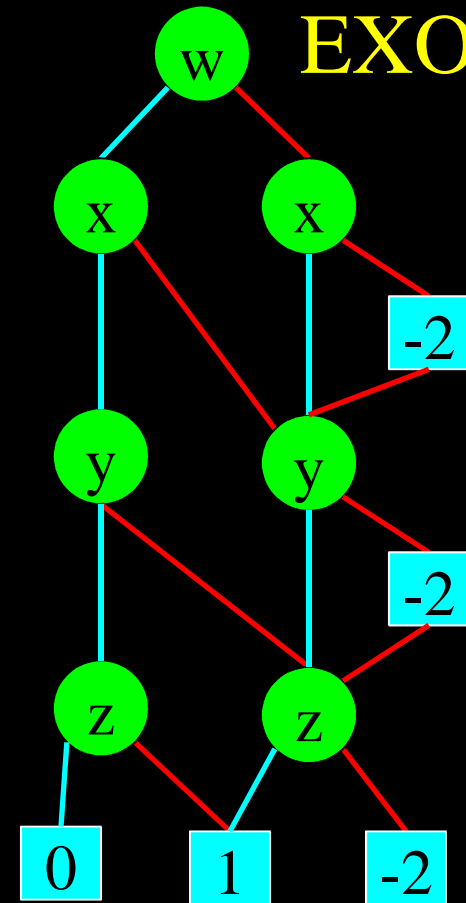
AND



OR



EXOR



Kronecker Multiplicative (K^*)BMDs

- An edge has two integer (rational) weights
- The weights are combined **additively and multiplicatively**
- The value of the function is derived by following a path from the root to a leaf and **adding/multiplying** edge weights encountered
- Three types of expansion are used: **Shannon, Positive Davio, Negative Davio**
- **Normalization rules** are quite complicated

Expansion Rules

- Shannon Expansion

$$\langle (a,m), F \rangle = a + m((1-x)^*F0 + x^*F1)$$

- Positive Davio Expansion

$$\langle (a,m), F \rangle = a + m(F0 + x^*F1)$$

- Negative Davio Expansion

$$\langle (a,m), F \rangle = a + m(F0 + (1-x)^*F1)$$

K*BMD Normalization Rules

- There is only one leaf labeled 0
- The low-edge of a node always has additive weight 0 and multiplicative weight 1
- The multiplicative weight at the high-edge is 1, if $F_0 = 0$. If the high-edge points to the leaf, then the multiplicative weight is normalized to 1

*BMD vs. K*BMD for Integer Encoding

